### Green Scheduling for Energy-Efficient Building Controls



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### Green Scheduling for Energy-Efficient Building Controls



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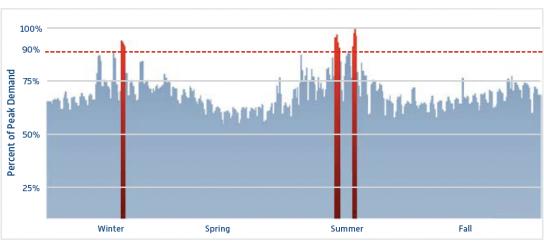
# The Peak Power Minimization Problem

supply = demand

#### 111 million viewers watch the SuperBowl

During a commercial, millions of refrigerators and microwaves trigger simultaneously, causing massive spikes in the energy demand

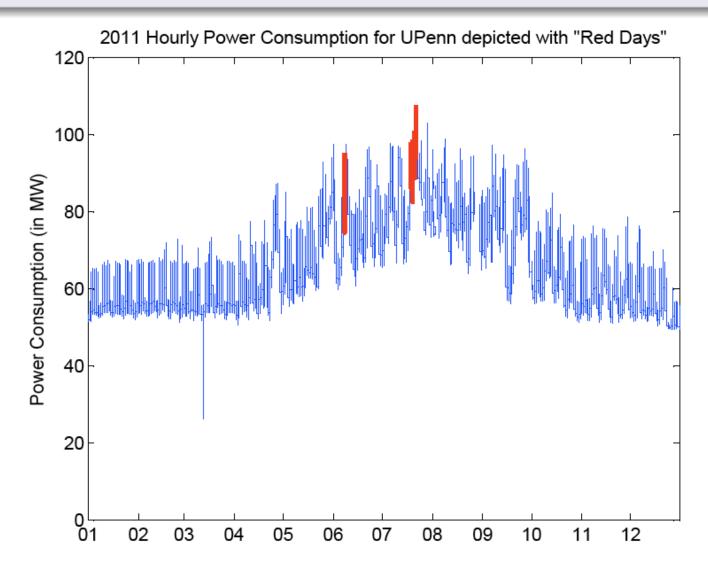




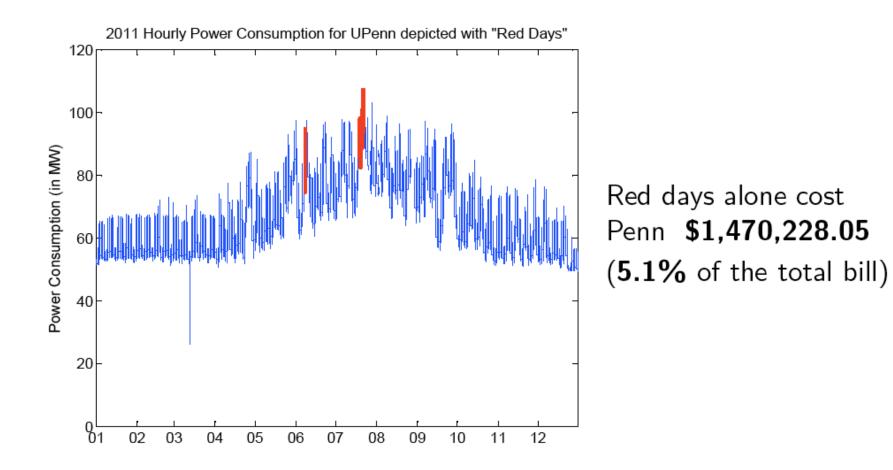
Human behavior and environmental conditions are responsible for high temporal correlation of energy demand.

#### Penn's Power Consumption

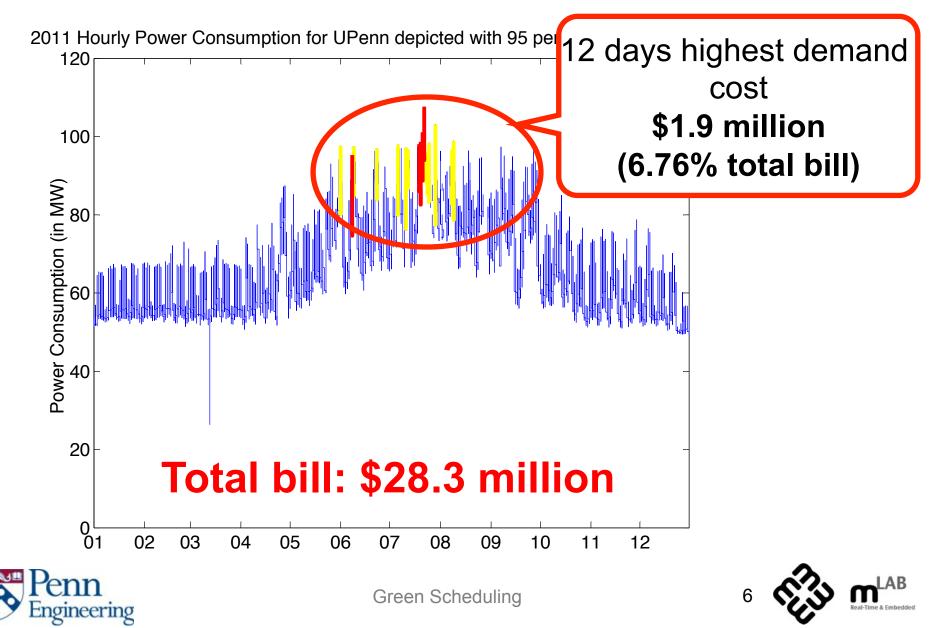
#### Estimated Annual Electricity Bill for UPenn is \$28,330,027.75



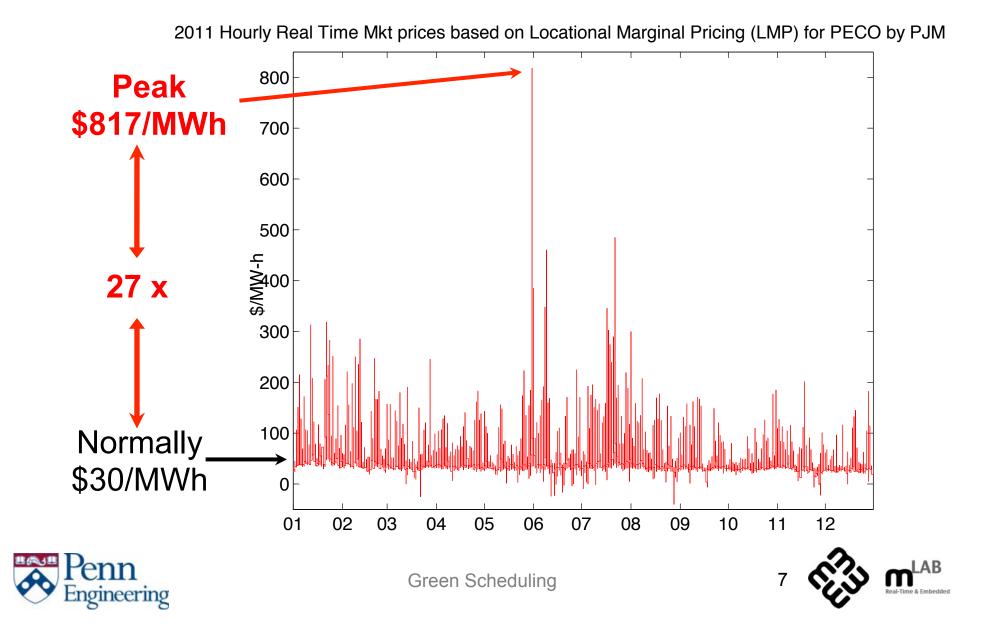
Demand across the grid (and hence the price of electricity) is very high



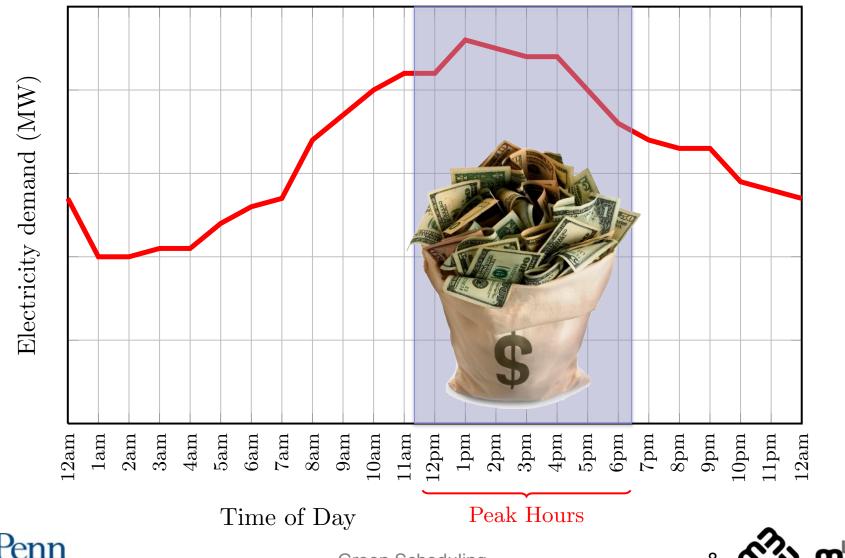
# Penn's Electricity Demand in 2011



### **Peak Electricity Demand is Expensive!**



### **Peak Electricity Demand is Expensive!**

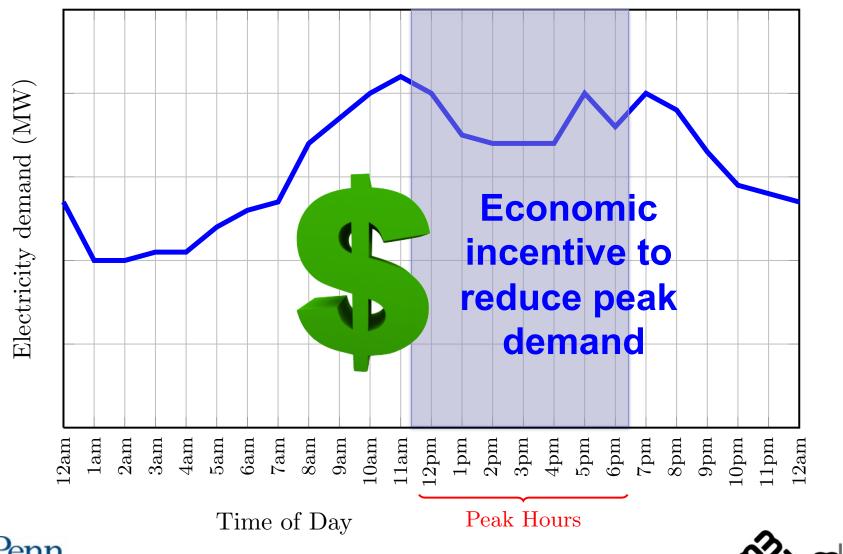




Green Scheduling



## **Peak Demand Reduction**

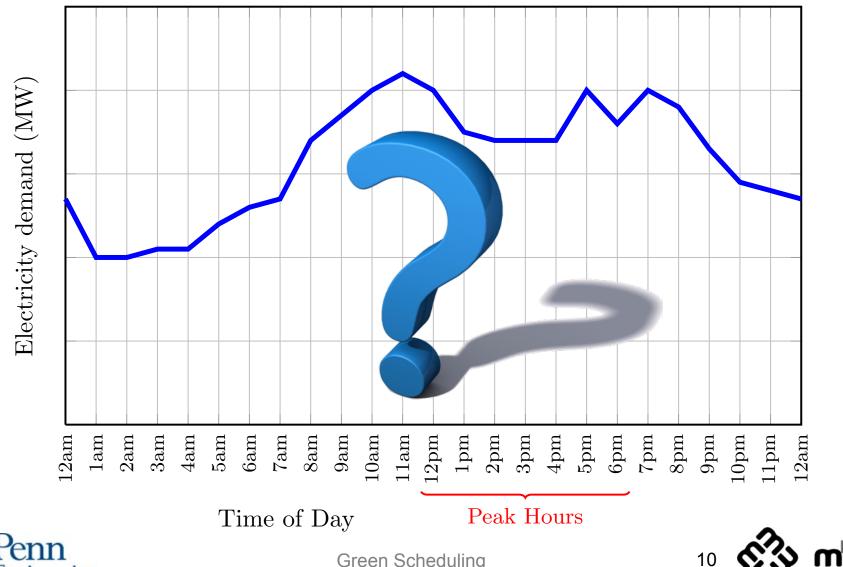




Green Scheduling

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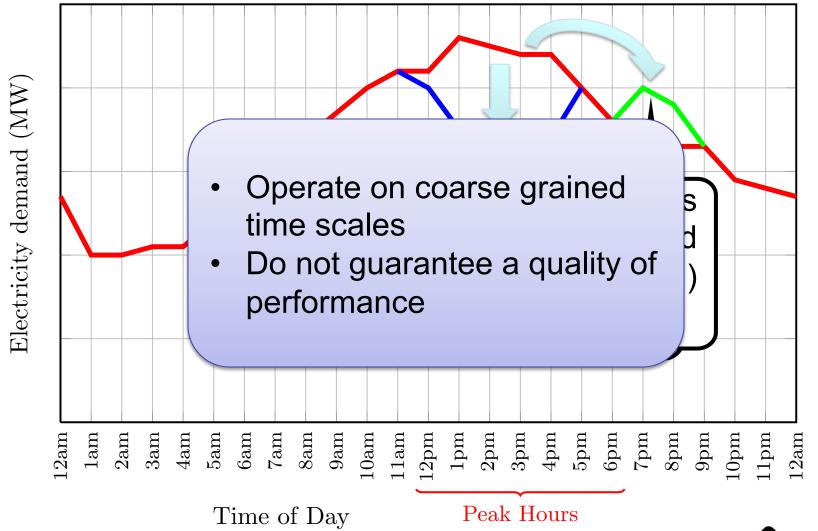
## **How to Reduce Peak Demand?**





Green Scheduling

### **Peak Demand Reduction Approaches**





Green Scheduling



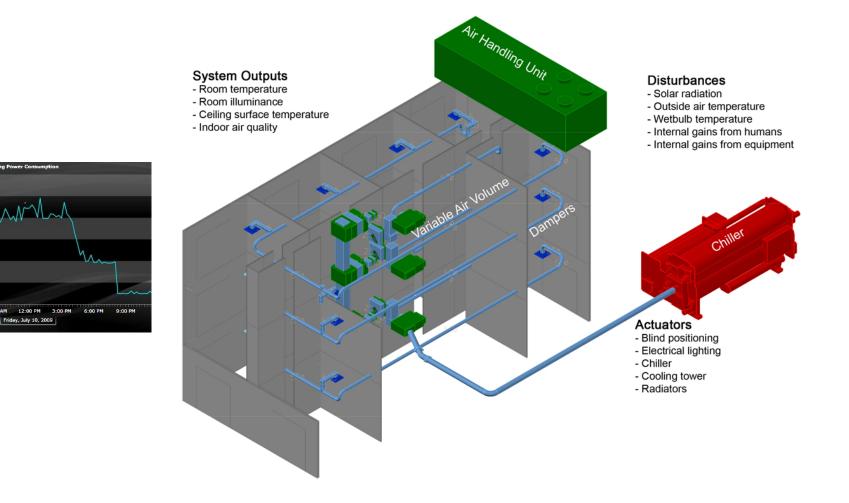
### Motivation

9:00 AM

3:00 AM

### **Un-coordinated Control Systems**

- HVAC (Heating, Ventilation and Air Conditioning) systems, chiller systems and lighting systems operate independently of each other
  - frequently result in temporally correlated energy demand surges (peaks)



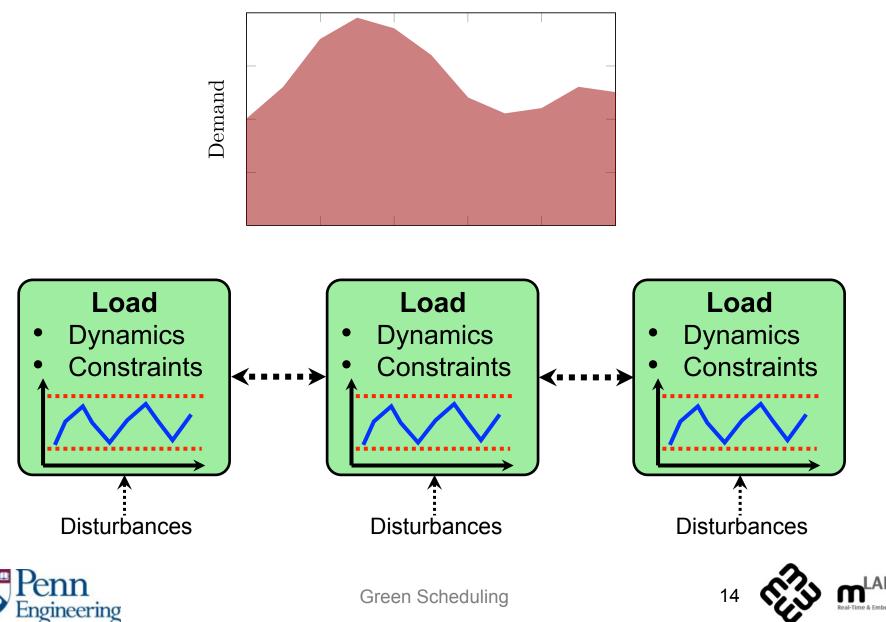
# **GREEN SCHEDULING APPROACH**



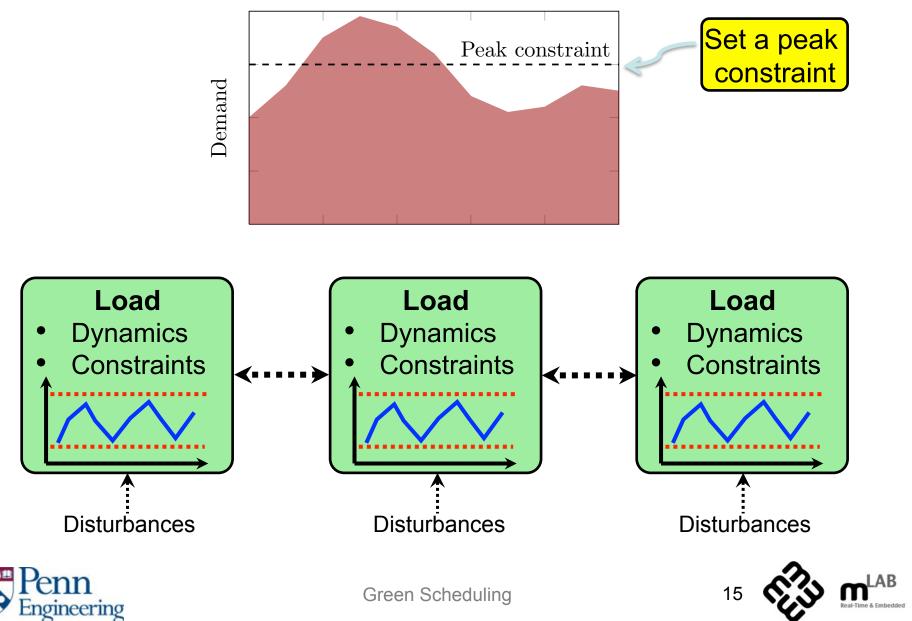
Green Scheduling



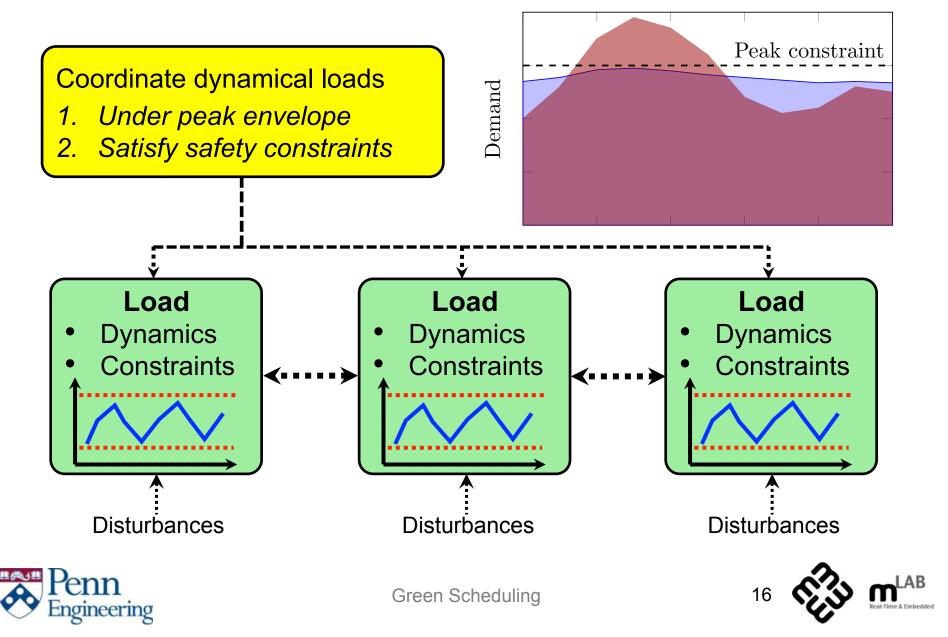
## **Green Scheduling (GS) Approach**



# **Green Scheduling (GS) Approach**

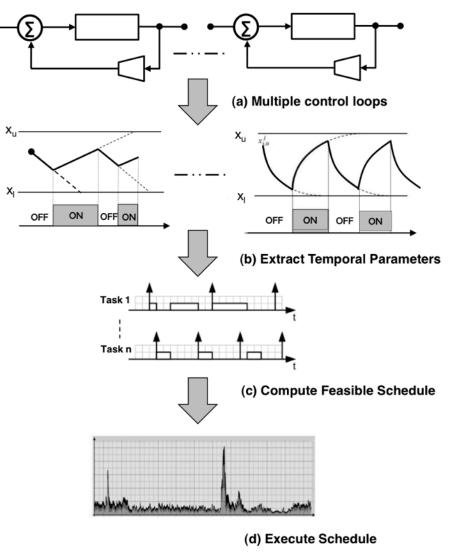


# **Green Scheduling (GS) Approach**



# **From Control to Scheduling**

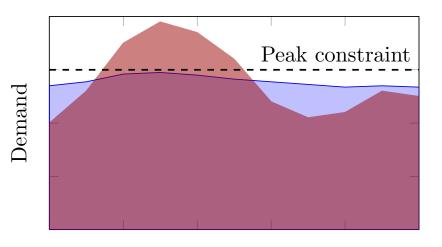
- Control loops are abstracted as tasks
- Extract temporal parameters across multiple control loops
- Compute a global schedule, reduce peak power by de-correlating systems





17 Keal-Time & Embedded

## **GS: Analysis & Synthesis**



Coordinate dynamical loads

- 1. Under peak envelope
- 2. Satisfy safety constraints

#### Schedulability analysis

Is a peak constraint feasible? (how to choose a peak constraint?)

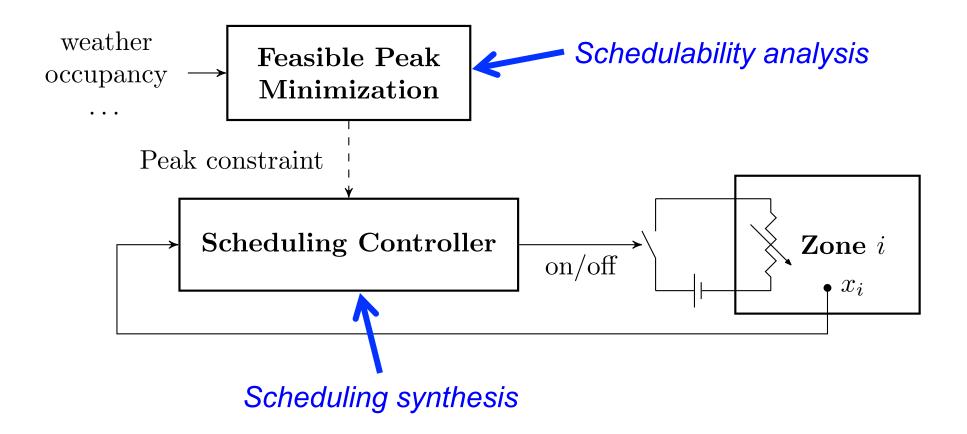
How to schedule the loads safely under a feasible peak constraint?

#### Schedule/Control Synthesis





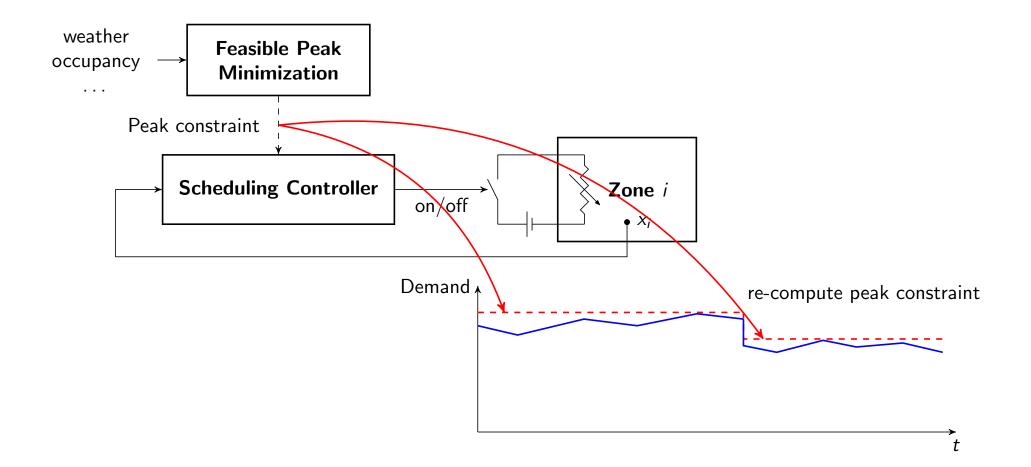
# **GS Control/Scheduling Structure**







# **GS Control/Scheduling Structure**







### From Control to Scheduling

#### In Energy Control Systems

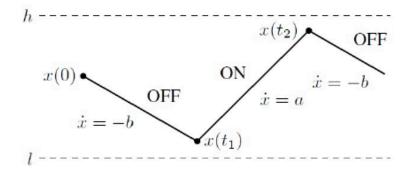
- Execution time is dependent on
  - Plant Dynamics: dimensions of the room, ingress and egress airflow
  - Environmental Conditions: outside weather, human occupancy, air quality
  - Initial State
- Tasks have **elastic execution times** where a task may have to perform more work, the longer its response time.
- Aim: keep the state of a system within a deadband

#### Both are resource contrained problems

- Here the resource is electricity/energy as opposed to a processor
- PSU (power supply unit) scheduling instead of CPU scheduling.

#### Task Model

- Tasks satisfy this differential equation:
  x(t+1) = dyn(m(t), x(t))
- Two operational modes,
  M = {ON, OFF}



• Linear tasks have dynamics defined as:

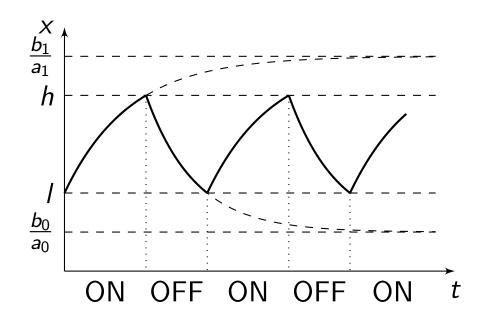
$$\dot{x}(t) = dyn(m(t), x(t)) = \begin{cases} a & \text{if } m(t) = \text{ON} \\ -b & \text{if } m(t) = \text{OFF} \end{cases}$$

 A task set *T* is *schedulable* by a policy *π* if *π* can schedule the tasks in *T* so that they are all safe and that all system-wide constraints are met.

#### **De-correlation** Constraint

At most one task can be in the ON mode at any time

#### First Order Task Model



Simplified heat balance equation to model each zone

$$C_i \frac{\mathrm{d}x_i}{\mathrm{d}t} = K_i \left( T_a - x_i \right) + Q_i \tag{1}$$

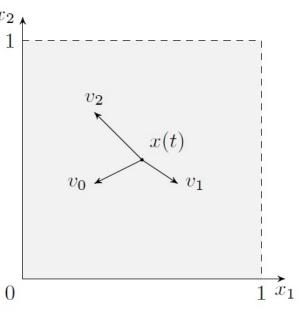
Zone temperature is governed by the following affine differential equation:

$$\frac{\mathrm{d}x_i}{\mathrm{d}t} = -\frac{K_i}{C_i}x_i + \left(\frac{K_i}{C_i}T_a + \frac{Q_i}{C_i}\right) = -a_ix_i + b_i \tag{2}$$

### Geometric Interpretation

Intuitive and simple framework for scheduling a system of linear tasks.

- Two linear tasks  $T_1$  and  $T_2$ , normalized so that their bounds are both [0, 1]
- Define a 2-dimensional state vector  $x = [x_1, x_2]^T \in \mathbb{R}^2$



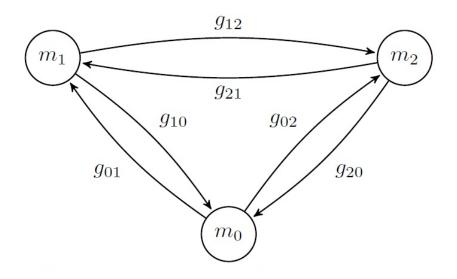
There are three scheduling modes:

- Mode 0:  $T_1$  and  $T_2$  are OFF (vector  $v_0$ )
- Mode 1:  $T_1$  is ON and  $T_2$  is OFF  $(v_1)$
- Mode 2:  $T_1$  is OFF and  $T_2$  is ON  $(v_2)$

#### Scheduling Policy

Keeps x(t) within bounds (invariant set) using mode vectors  $v_0$ ,  $v_1$  and  $v_2$ 

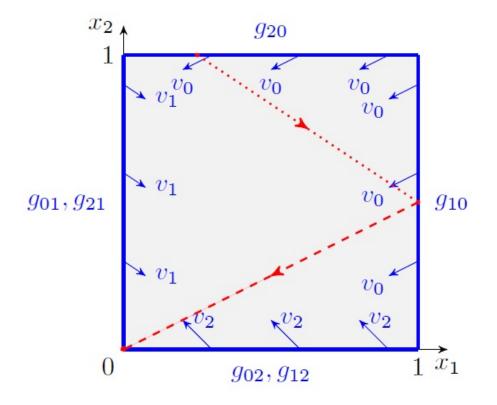
### Scheduling Policy as Hybrid Automaton



- States  $\stackrel{correspond to}{\rightarrow}$  scheduling modes
- Edges  $\stackrel{correspond to}{\rightarrow}$  switching between modes
- $g_{ij}$  is the guard associated for each edge, for the transition from mode *i* to mode *j*.
- Scheduling policy  $\pi$  for the task set is simply a set of guards  $\{g_{ij}\}$

#### Lazy Scheduling Policy: Hybrid Automaton

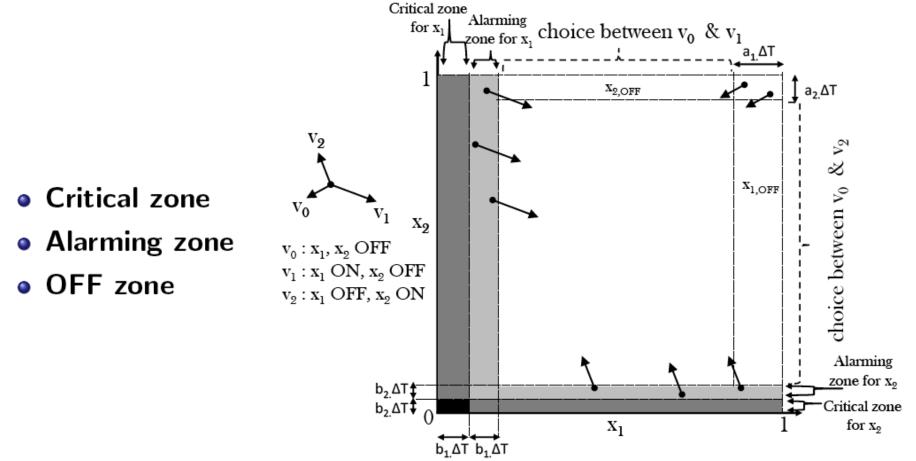
Lazy Policy: All tasks stay in their current modes as long as they are safe



- $g_{01}$  and  $g_{21}$  are both  $(x_1 \leq l_1)$ ;
- $g_{02}$  and  $g_{12}$  are both  $(x_2 \leq l_2)$ ;
- $g_{10}$  is  $(x_1 \ge h_1 \land x_2 > l_2)$ ;
- $g_{20}$  is  $(x_2 \ge h_2 \land x_1 > l_1)$ .

### The Lazy Scheduling Algorithm

*Lazy*, because switching decisions are made only if a task is approaching its thresholds



### Simulation: Two-Task System

- Feasibility constraint: Keep temperature centered around mean 70°F
- ${\bullet}\,$  Heating system operates with a power of 12000 BTU/h or 3.517 kW
- Cooling occurs through heat loss and does not consume any extra power
- Each time step of the algorithm is of 15 minute duration

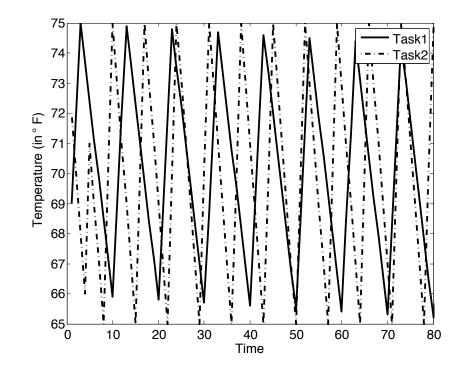


Figure: Peaks occur when tasks run independently

### Simulation: Peak Reduction

#### Peak Reduction of 50%

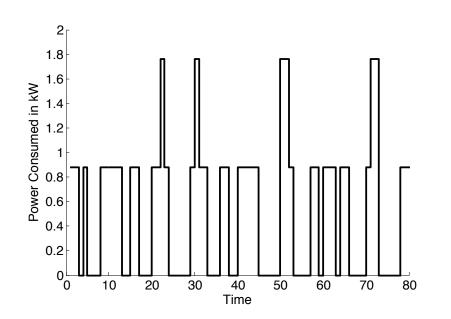


Figure: Power Consumption: Independent Tasks

- Peak Power = 1.758 kW
- Total Energy = 50.11 kWh

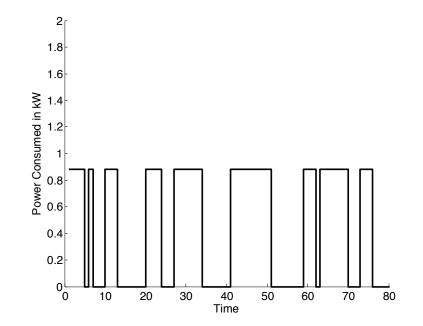


Figure: Power Consumption: Lazy Scheduling

- Peak Power = 0.879 kW
- Total Energy = 45.72 kWh

### Simulation: Lazy Scheduling

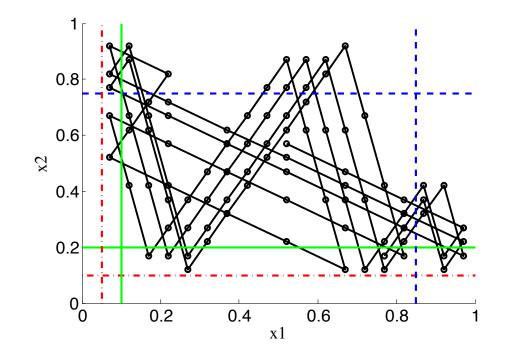


Figure: State-Space Trajectory

- Tasks remain within thresholds
- Modes only change in a *lazy* manner (near thresholds)

### Simulation: Lazy Scheduling for 3 Tasks

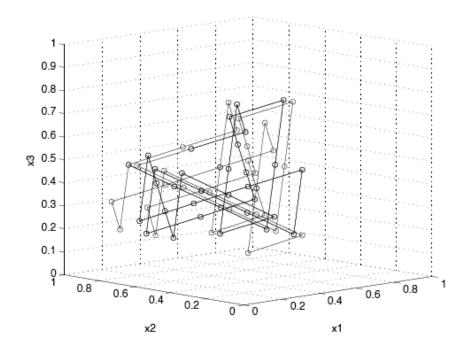
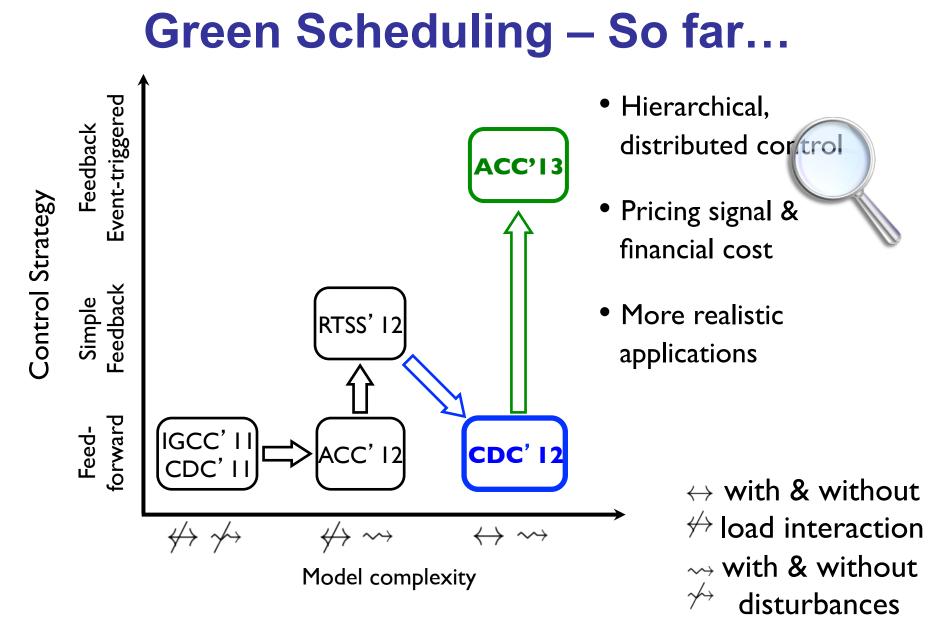


Figure: State-Space Trajectory





Engineering

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### **APPLICATION:** ENERGY-EFFICIENT OPERATION OF MULTIPLE CHILLER PLANTS



Green Scheduling



## **Peak Demand Reduction**

Chiller

Supply

Side

Building

Demand Side



- ✓ Chillers
- ✓ Cooling Towers
- ✓ Thermal Energy Storage
- ✓ Condensers
- ✓ Underground mains ...



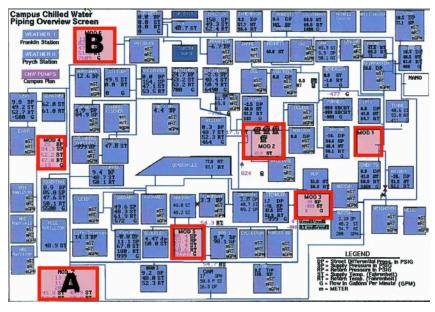
- ✓ Variable Air Volume Box (VAV)
- ✓ Air Handling Units
- ✓ Occupant comfort
- ✓ Solar heat gain ...



Green Scheduling



## **Chiller Plant Operation at Penn**





Chilled water distribution system at the University of Pennsylvania Chiller Plant A

- Over 4M gallons of chilled water (42 F) pumped into the campus.
- Plant A consumes **26 MW** at full capacity
- Plant A and Plant B account for > 30% of total peak power consumption (108 MW)





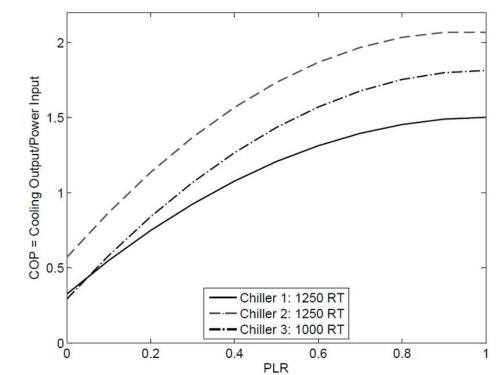
## COP vs. PLR

Approximate COP by quadratic function of PLR

$$COP = a_0 + a_1 PLR + a_2 PLR^2$$

#### Example:

If cooling load = 1000 RT<sup>1</sup>, chiller capacity = 1250 RT, then PLR = 1000/1250 = 0.8, and COP obtained from COP-PLR curve.



<sup>1</sup> Refrigeration Ton is the amount of heat that must be removed to melt 1 Ton of ice in 24h. 1 RT = 3517W

## **Thermal Energy Storage (TES)**

Thermal Energy Storage is used for demand shifting and off loading chillers.

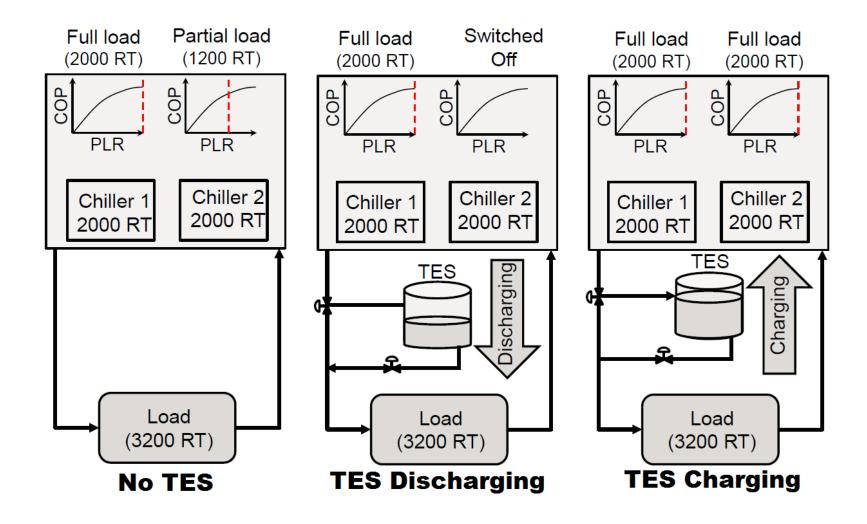
- Long term: > 10 hours
- **Short term**: < 2 hours



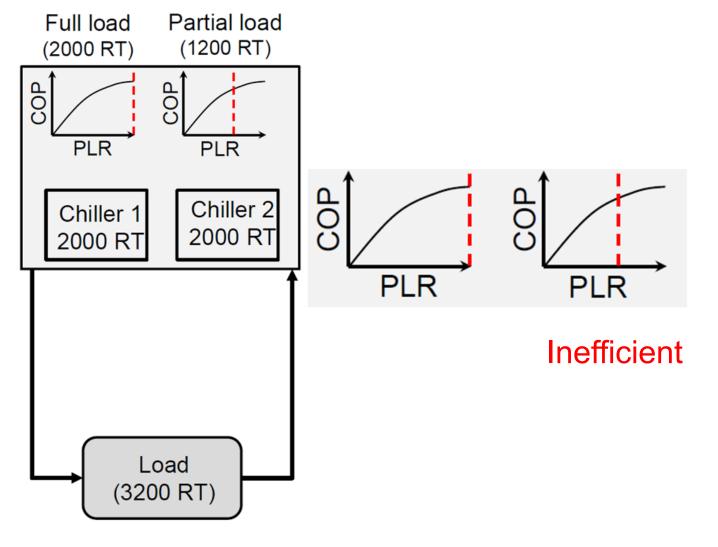
TES at Cornell University

Short-term Thermal Energy Storage can improve the COP of the chiller plant.

## How TES Improves the COP?



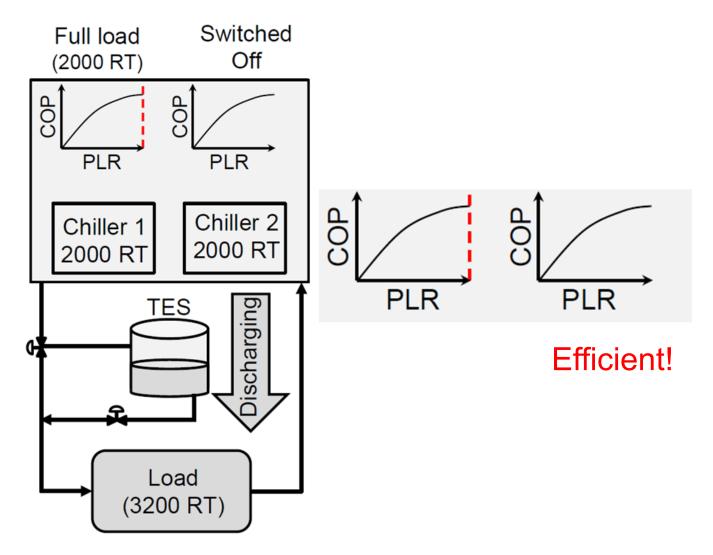
## **No Thermal Energy Storage**







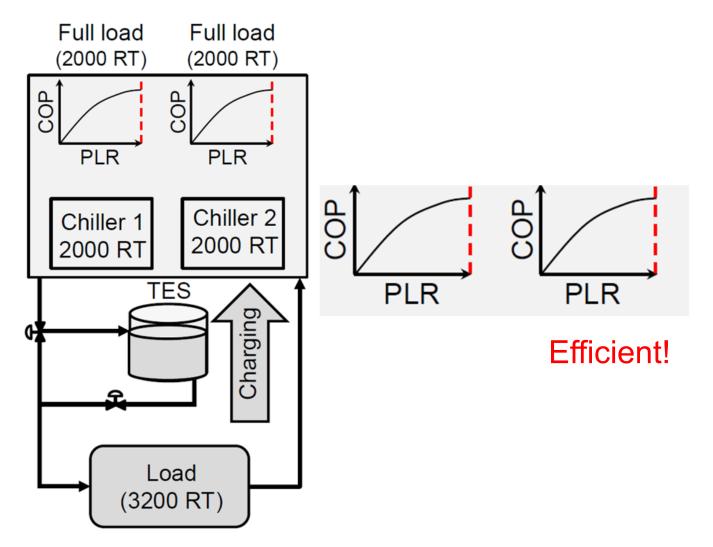
## **TES Discharging**







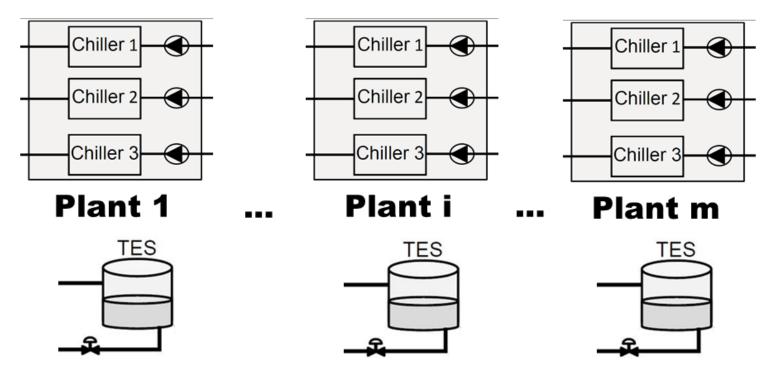
## **TES Charging**







## Multiple chiller plants with TES



*Uncoordinated operation* among multiple chiller plants can cause large spikes in total electricity consumption.

Our goal is to coordinate multiple chiller plants to reduce their aggregate peak power consumption.

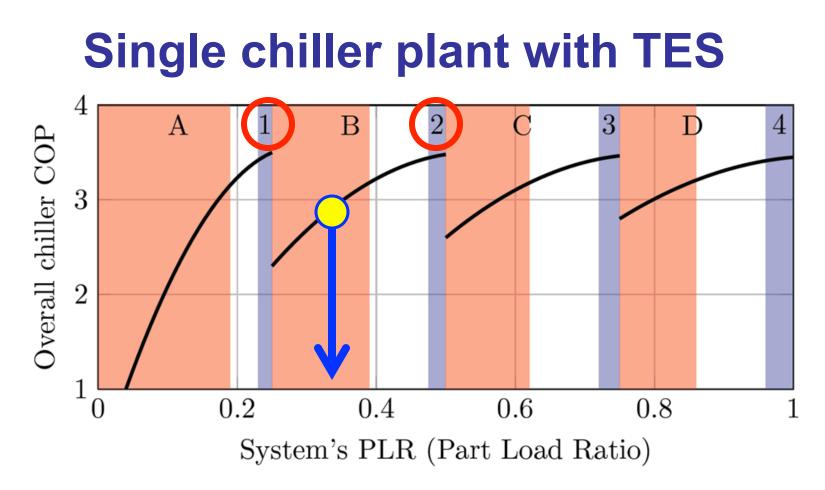
## **GS for Chiller Plants**

- Consider m > 1 chiller plants.
- Each plant has a short term TES system.
- Compute modes for each hour h, h = 0,1,...,H, based on load forecasting.
- Charging mode: water level increases with rate a<sub>i,h</sub> > 0;
- Discharging mode: water level decreases with rate b<sub>i,h</sub> < 0.</li>

Behl, M.; Nghiem, T. X.; Mangharam, R., *Green Scheduling for Energy-Efficient Operation of Multiple Chiller Plants*, RTSS 2012.







- Say PLR = 0.34 i.e., in region B
- Plant can operate in optimal regions 1 and 2 if a TES is available.

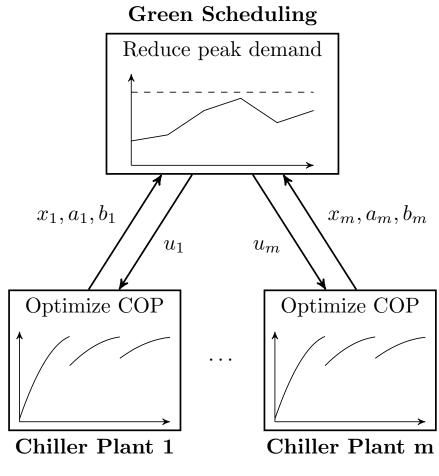




## **GS for Chiller Plants**

Schedule the operating modes  $u_i(t)$ ,  $t \ge 0$ , of all plants such that: Green Scheduling

- Safety Constraint:  $I_i \leq x_i(t) \leq h_i \quad \forall t, i$
- Peak Constraint:  $\sum_{i=1}^{n} u_i(t) \le k(t) \quad \forall t$



## **Simulation Setup**

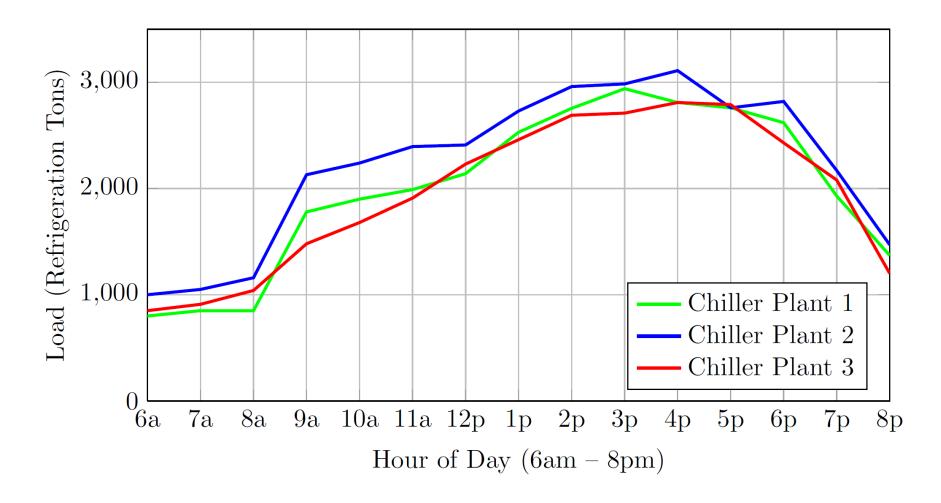
Three chiller plants, each containing 3 chillers.

- Lower safety threshold: 3m
- Upper safety threshold: 13.5m
- Time step = 15 mins

#### Chiller plant configuration in the simulation

Plant 1	3 chillers rated at 1250 RT, 1200 hp
Plant 2	3 chillers rated at 1250 RT, 1200 hp
Plant 3	3 chillers rated at 1000 RT, 900 hp
T <sub>chws</sub>	5.5°C (42°F)
T <sub>cwr</sub>	20°C (68°F)
$\Delta T$	10°C

## **Average Hourly Load Profile**







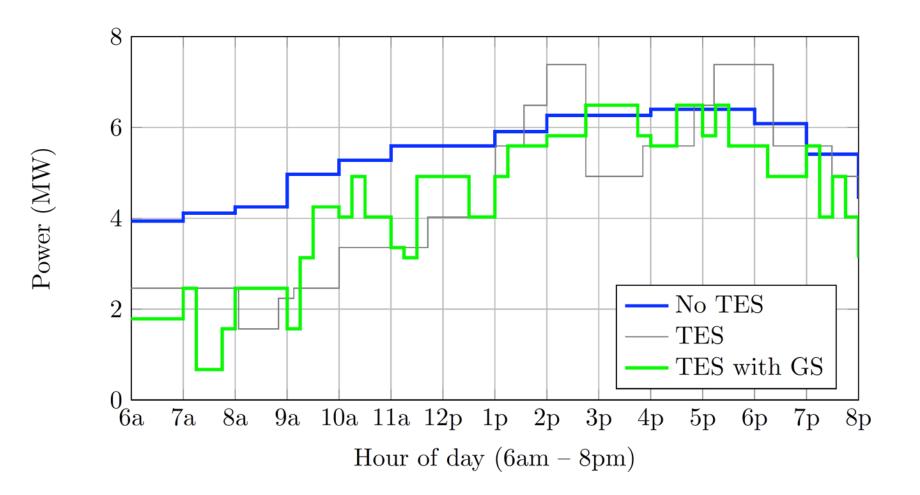
## **Simulation Scenarios**

- 1. Case 1: No Thermal Energy Storage present
- 2. Case 2: Thermal Energy Storage with uncoordinated operation
- **3. Case 3:** Thermal Energy Storage with (discrete-time) Green Scheduling





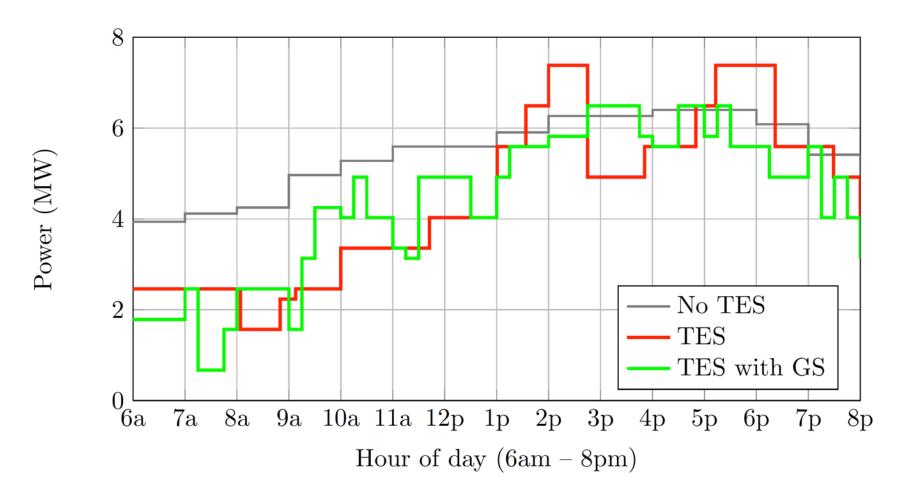
## **Simulation Result: Power Demand**







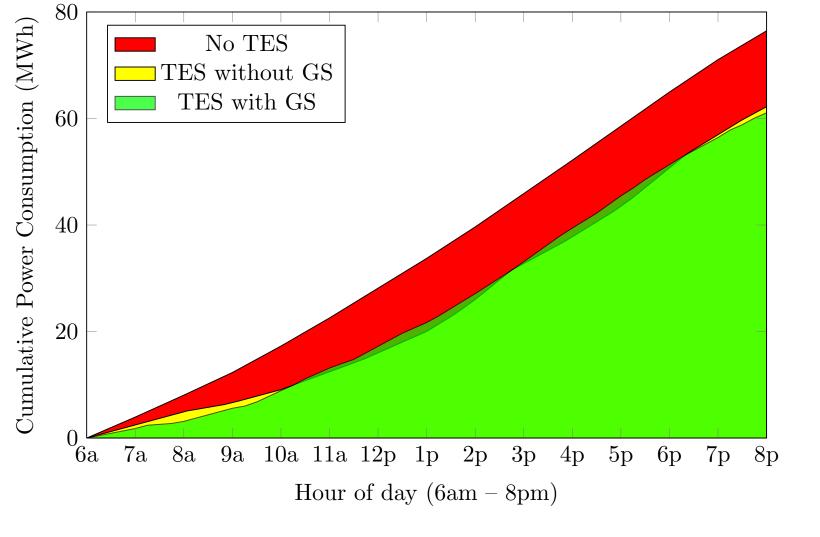
## **Simulation Result: Power Demand**







## **Cumulative Energy Consumption**







## **Electricity Pricing Policy**

PECO's demand pricing rate structure:

Block	kWh's in Block	Charges (cents per kWh)
First Block	80 × peak	24.94
Second Block	80 × peak	12.67
Third Block	Remaining	8.64





## **Simulation Result Summary**

#### Green Scheduling leads to the lowest electricity bill.

	Peak (MW)	Energy Consumption (MWh) single day	Expected Monthly Bill (\$)	% savings
No TES	6.40	80.9	292,801	-
On-Off with TES	7.38	66.1	274,266	6.33
GS with TES	6.48	61.4	243,461	16.85





## **SCHEDULING SYNTHESIS:** *DISCRETE-TIME FEEDBACK SCHEDULING*





## **Discrete-time Green Scheduling**

- Consider discrete-time dynamics  $x(t+1) = Ax(t) + B_0 + Bu(t) + Wd(t), \ \forall t \in \mathbb{N}$ for a finite time horizon [0, T].
- Safety constraint:  $x(t) \in S \subset R^n \ \forall t$
- Initial state  $x(0) \in S$
- $X_f \subseteq S$ : set of desired final states at time T:  $x(T) \in X_f$
- Define the set of **admissible** discrete-time schedules

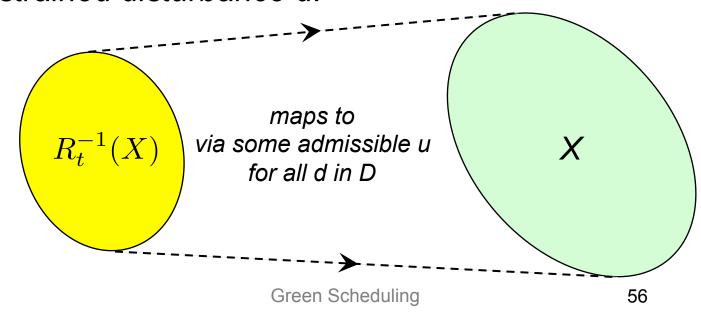
Define  $\mathbb{U}(k(\cdot), S, x_0, X_f)$  as the set of all discrete-time schedules  $u: \{0, 1, \dots, T-1\} \rightarrow \{0, 1\}^m$  such that: 1  $\sum_{i=1}^m u_i(\tau) \le k(\tau)$  for all  $\tau$ ;  $\leftarrow$  Peak constraint 2  $x(\tau) \in S$  for all  $\tau$ ;  $\leftarrow$  Safety constraint 3  $x(T) \in X_f$ .  $\leftarrow$  Terminal state is reached

### **Robust Backward Reachable Sets**

Robust backward reachable set operator: Given any  $X \subseteq S$  at time step t, define operator

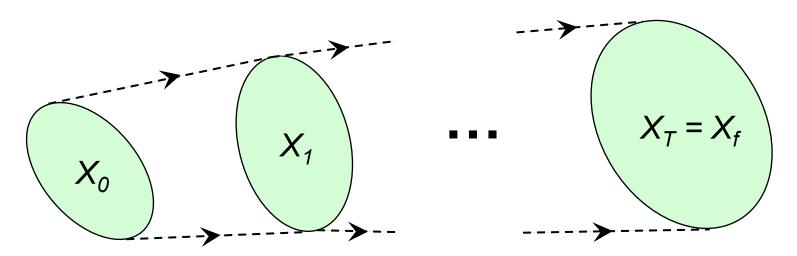
 $R_t^{-1}(X) = \{x \in S : \exists u : \sum_i u_i \le k(t) \land f(x, u, d) \in X \, \forall d \in D\}$ which is the set of safe states that can reach X with

some admissible control *u*, regardless of the constrained disturbance *d*.

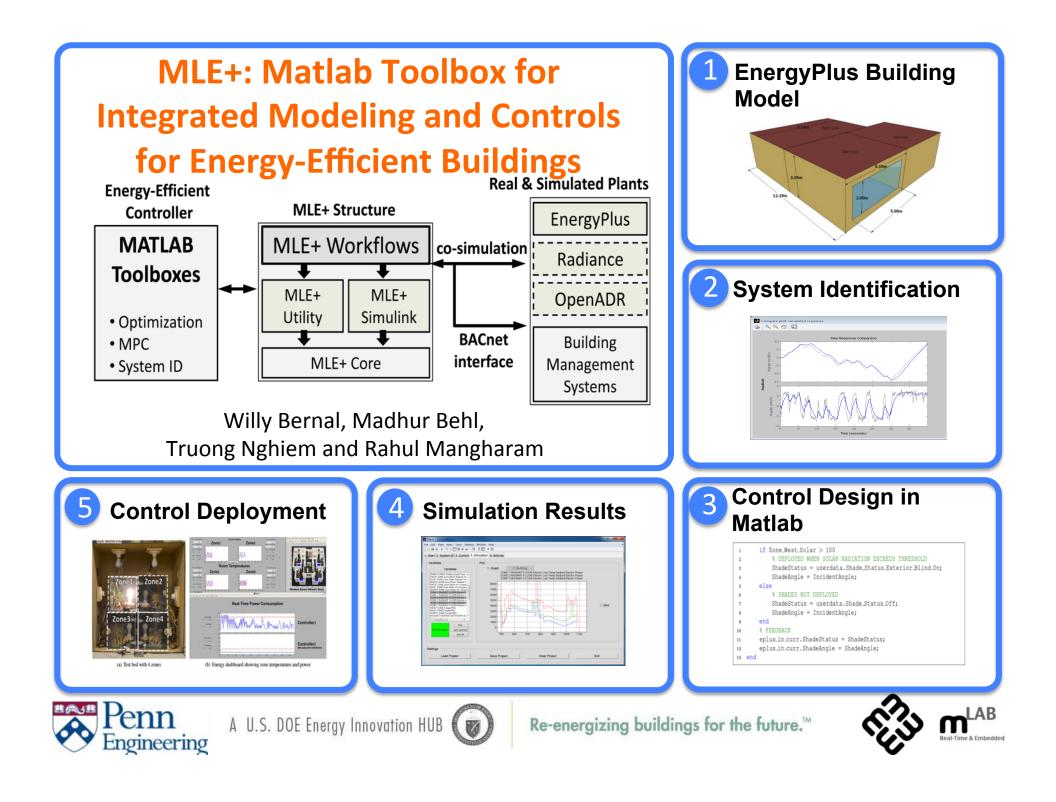


## **Pre-compute Backward in Time**

• Using robust backward reachable set operator, compute a sequence of sets { $X_t : t = 0, 1, ..., T$ } backward in time from  $X_T = X_f : X_t = R^{-1}_t(X_{t+1})$ 

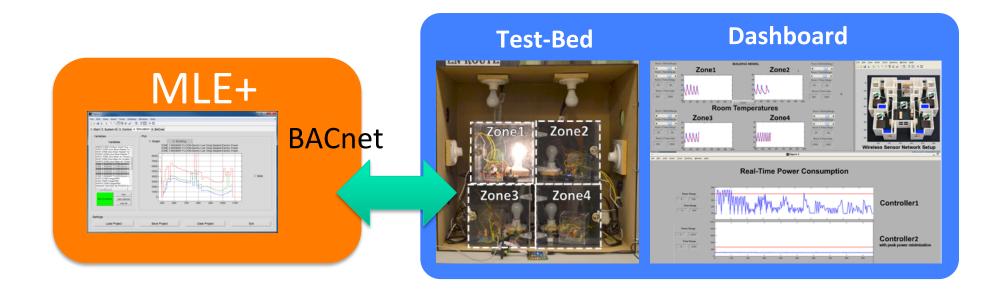


 Each X<sub>t</sub> is the set of safe states from which and from time step t, the system state can reach the final set X<sub>f</sub> safely under the time-varying peak constraint k(·).



### MLE+: Design and Deployment Integration of Energy-Efficient Building Controls

Willy Bernal, Madhur Behl, Truong Nghiem and Rahul Mangharam





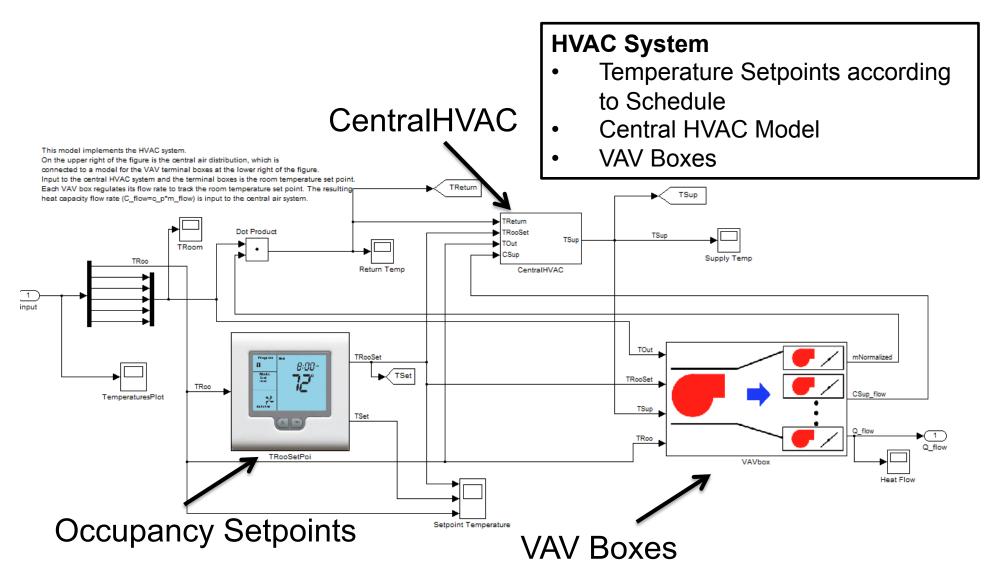
A U.S. DOE Energy Innovation HUB

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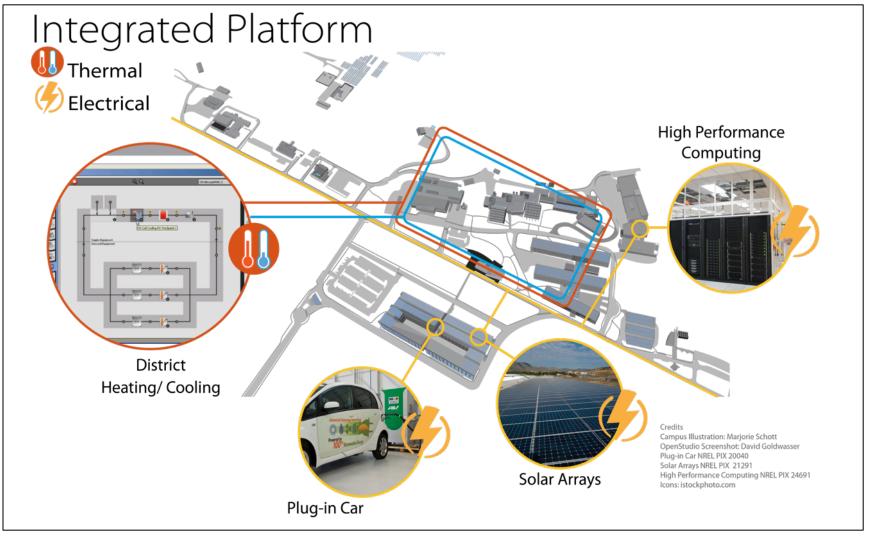
Re-energizing buildings for the future.™



### **Integrated Modeling: HVAC System**

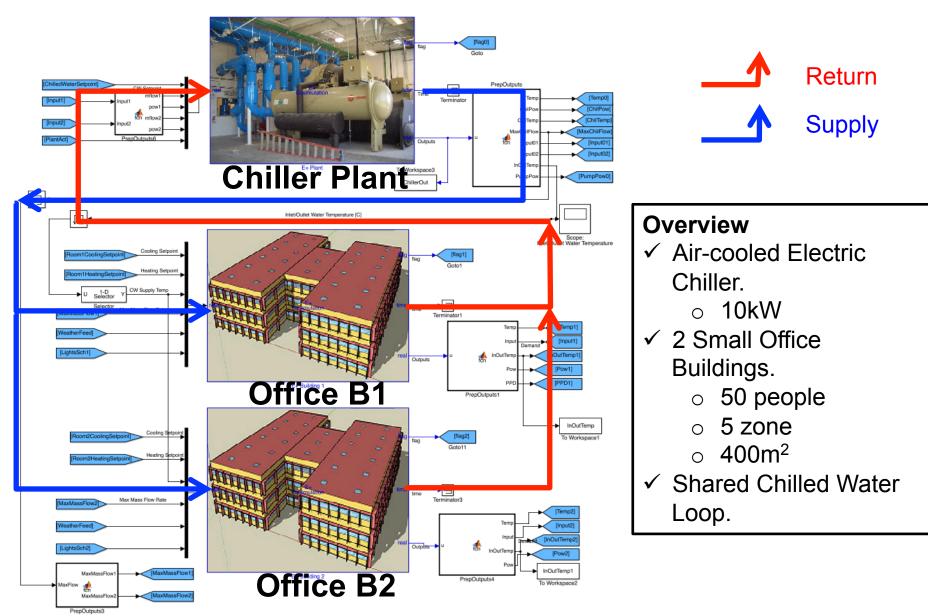


## **NREL: Campus-Wide Simulation**

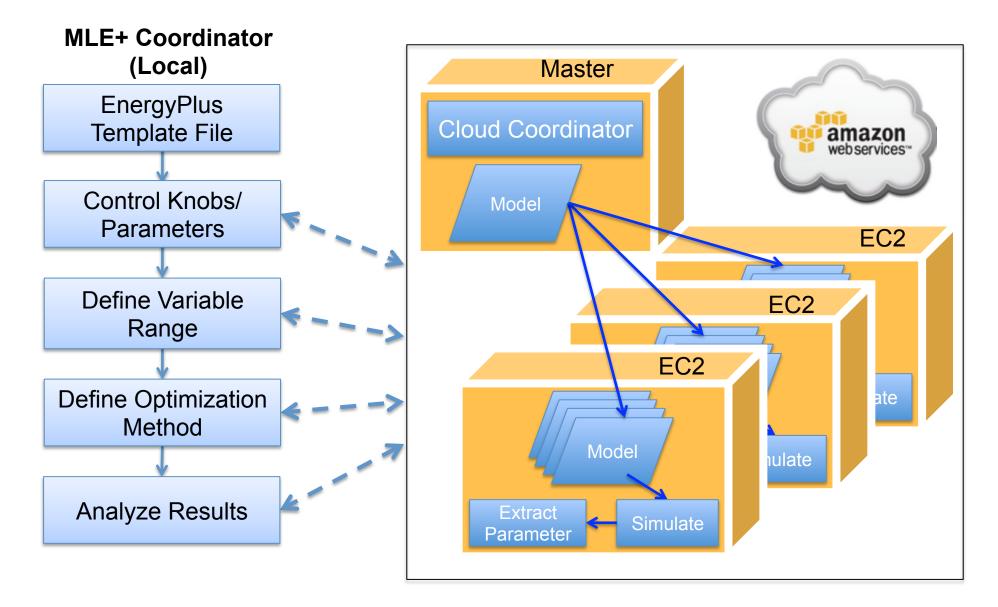


**Courtesy of NREL** 

#### **Campus Simulation: Two Building Example**



### **MLE+ Cloud Analytics: Overview**



## **Cyber-Physical Systems**

- Intersection of Computation, Controls & Communication
  - Safety-critical and Life-critical systems
- Tightly coupled with (Messy) Physical Plants
  - Interesting Domains: Medical, Energy, Automotive...
- New Interesting Problems involving:
  - Scheduling and Control
  - From Verified Models to Verified Code of Closed-loop Systems





# **Networked Cyber-Physical Systems**

Medical Device Software & Systems Network CPS Industrial Control Nets Automotive CPS Automotive Plug-n-Play

Real-Time Parallel Computing

Distributed Real-Time Systems and Real-Time Network Protocols

Modeling & Tools for Virtualization and Deployment

**Domain-specific Platforms** 

Energy-efficient Building Automation Industrial Control & Actuation Networks

Feedback-based Medical Device Networks



