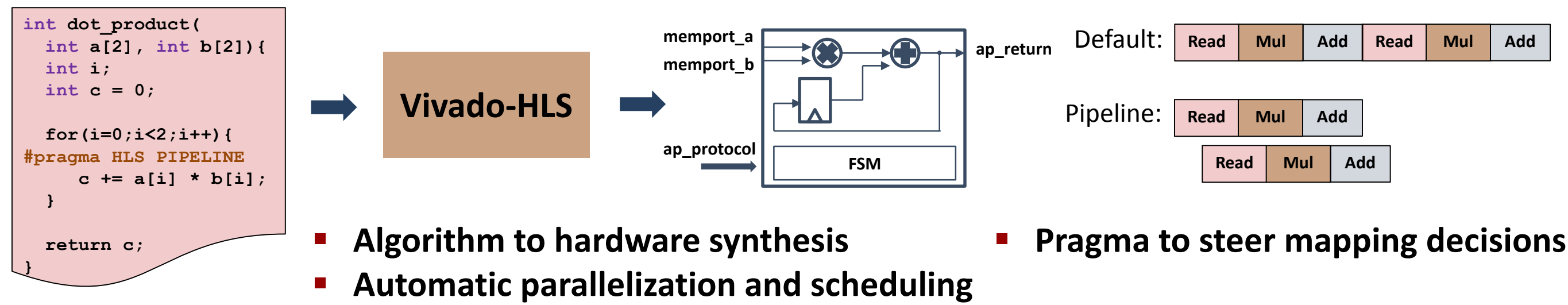


Using Vivado-HLS for Structural Design: a NoC Case Study

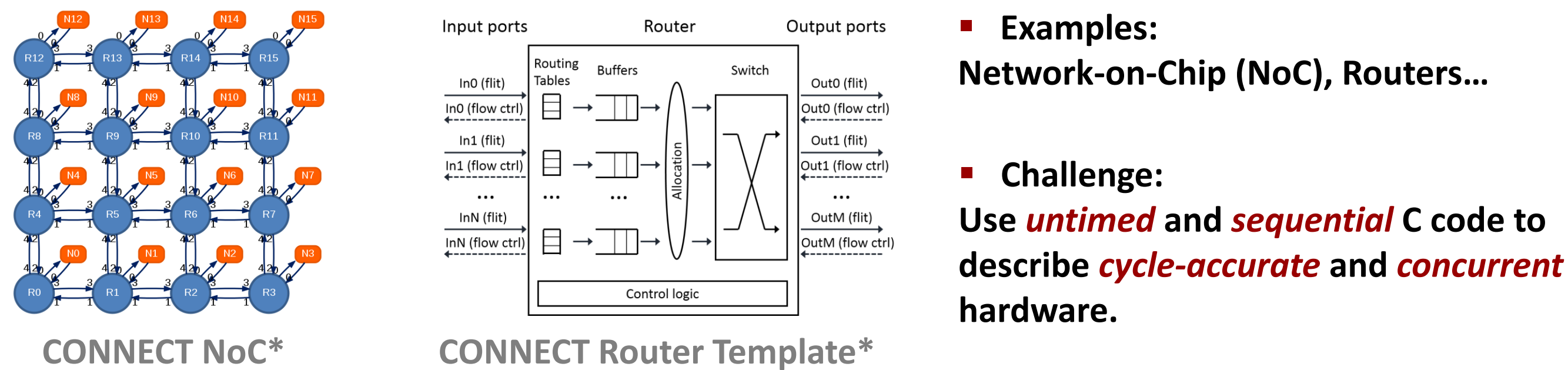
Zhipeng Zhao <zzhao1@andrew.cmu.edu>, James C. Hoe <jhoe@ece.cmu.edu>

Should we use HLS for structural design?

Xilinx Vivado High Level Synthesis (HLS)



Structural Design: *Precise Cycle- and Bit-level Control*



Can We and Should We Use Vivado-HLS for Structural Design?



Can we? Yes; Should we? Depends.

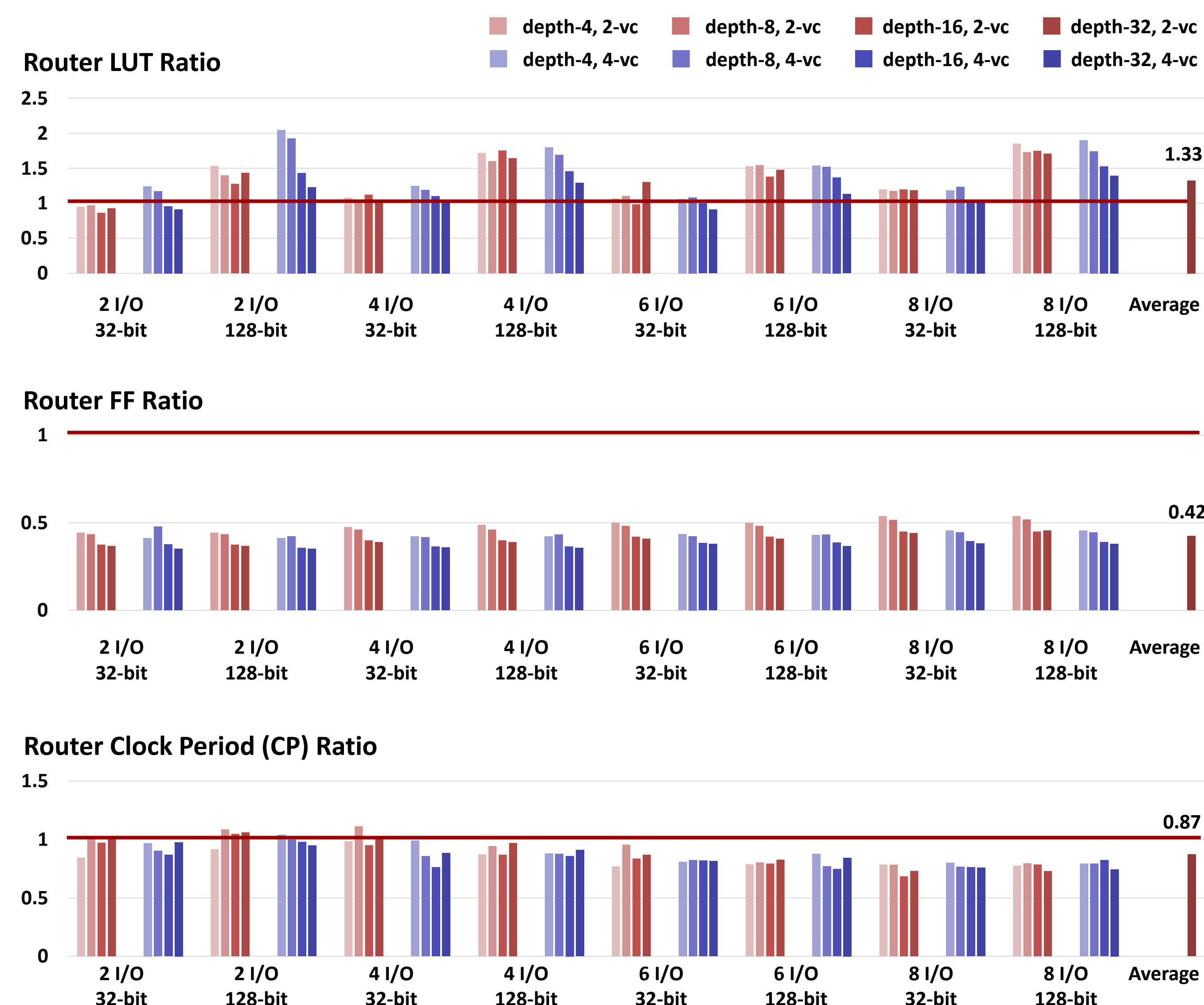
Case Study Overview

- RTL Reference: CONNECT Parameterized NoC (www.ece.cmu.edu/calcm/connect)
- Productivity gain from separation of *functionality* and *structural details*. But offer little advantage in pure netlisting
- Produced exact cycle- and bit-accurate replacements
- Achieved comparable hardware cost and critical path

Productivity Improvement

- Natural C's sequential reading is maintained. Utilize C's facilities to capture more maintainable and scalable designs. **1126** lines of C++ VS. **2605** lines of Bluespec System Verilog
- Still need to specify structural details to get desired hardware
- Separation of functionality and structural details enables fast design iteration
- Ordering discipline could be cumbersome when creating a netlist to compose submodules

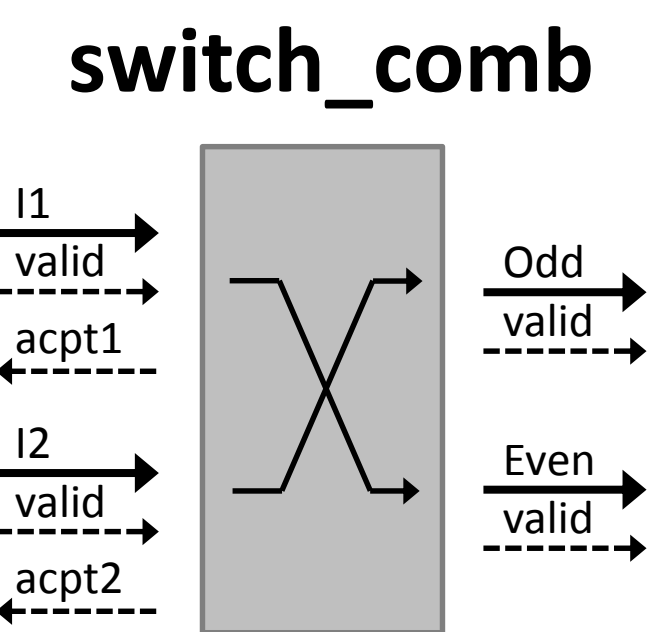
Hardware Quality (HLS VS. CONNECT)



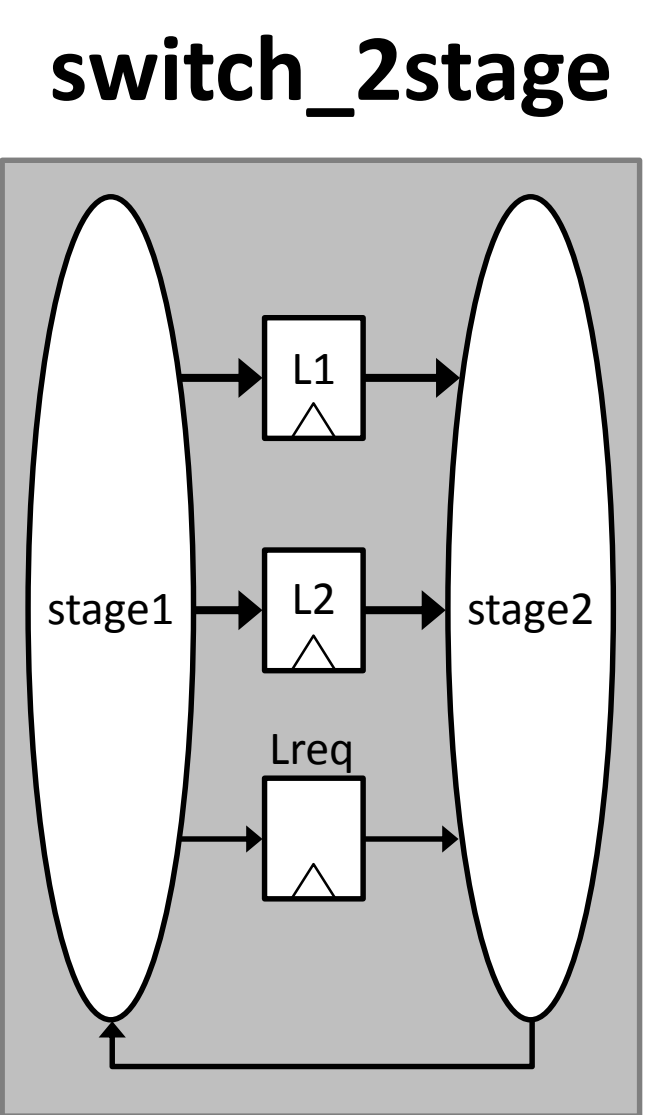
What worked very well

Illustrative example here; see website for complete router source

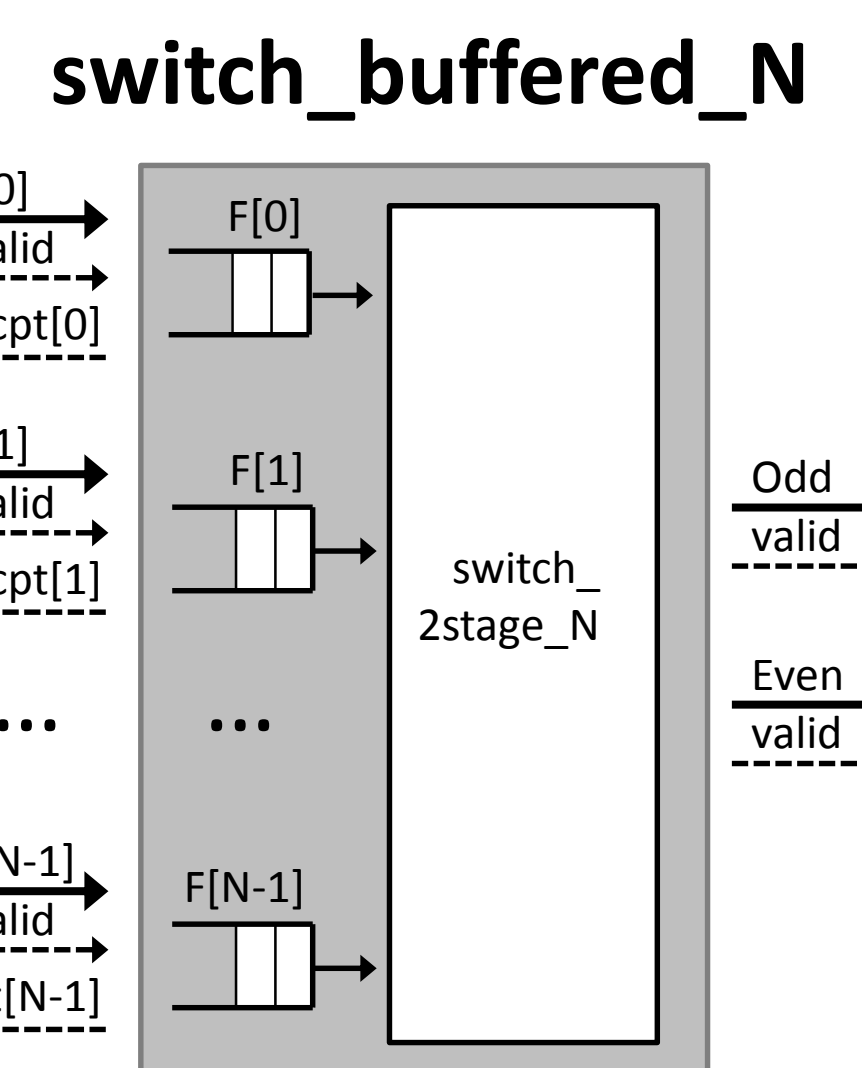
```
typedef struct vDat {bool v; int d;}
void switch_comb(vDat I1, vDat I2, vDat *Odd, vDat *Even,
                bool *acpt1, bool *acpt2) {
    bool *acpt1=(*Odd).v=(*Even).v=false;
    //For Odd
    if(I1.v && (I1.d%2))
        {(*Odd).v=true;(*Odd).d=I1.d;*acpt1=true;}
    else if (I2.v && (I2.d%2))
        {(*Odd).v=true;(*Odd).d=I2.d;*acpt2=true;}
    ... repeat for Even ...
}
```



```
void switch_2stage(vDat I1, vDat I2, vDat *Odd,
                  vDat *Even, bool *acpt1, bool *acpt2) {
    //Internal states
    static vData L1,L2;//Latch for I1 and I2
    static Path Lreq;//Latch for four possible connections
    /* --- stage 2 --- */
    Path grnt=allocate(Lreq);
    if (grnt.L1xOd)*Odd=L1;
    else if (grnt.L2xOd)*Odd=L2;
    else (*Odd).v=false;
    ... repeat for Even ...
    /* --- stage 1 --- */
    if (grnt.L1xOd|grnt.L1xEv)
        L1.v=false;
    if (I1.v && (!L1.v))
        {*acpt1=true; L1=I1;}
    else *acpt1=false;
    ... repeat for I2 and L2 ...
    Lreq=decode(L1,L2);
}
```



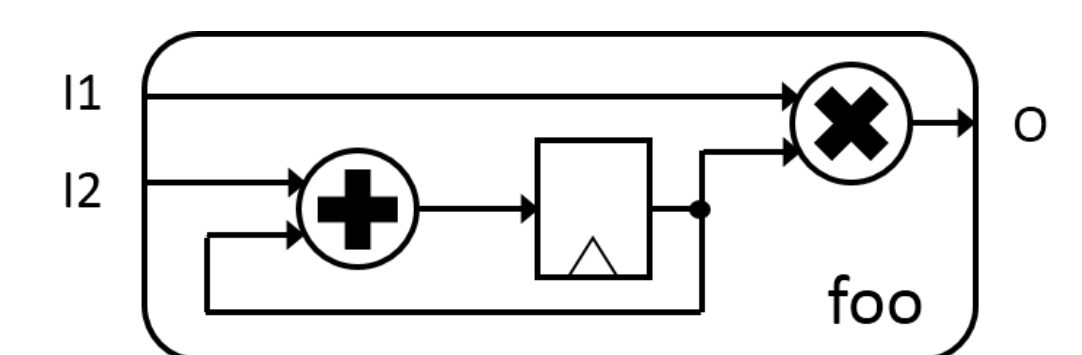
```
void switch_buffered_N(vData I[N], vData *Odd,
                      vData *Even, bool acpt[N]) {
    #pragma HLS ARRAY_PARTITION variable=acpt complete dim=1
    #pragma HLS ARRAY_PARTITION variable=I complete dim=1
    static FIFO<int> F[N];
    bool okX[N];
    vData frontX[N];
    for(int i=0;i<N;i++){
        #pragma HLS UNROLL
        frontX[i].v=!F[i].empty(); frontX[i].d=F[i].front();
    }
    switch_2stage_N(frontX, Odd, Even, okX);
    for(int i=0;i<N;i++){
        #pragma HLS UNROLL
        if(!F[i].full() && I[i].v)
            {F[i].push(I[i].d);acpt[i]=true;}
        else acpt[i]=false;
        if(okX[i]) F[i].pop();
    }
}
```



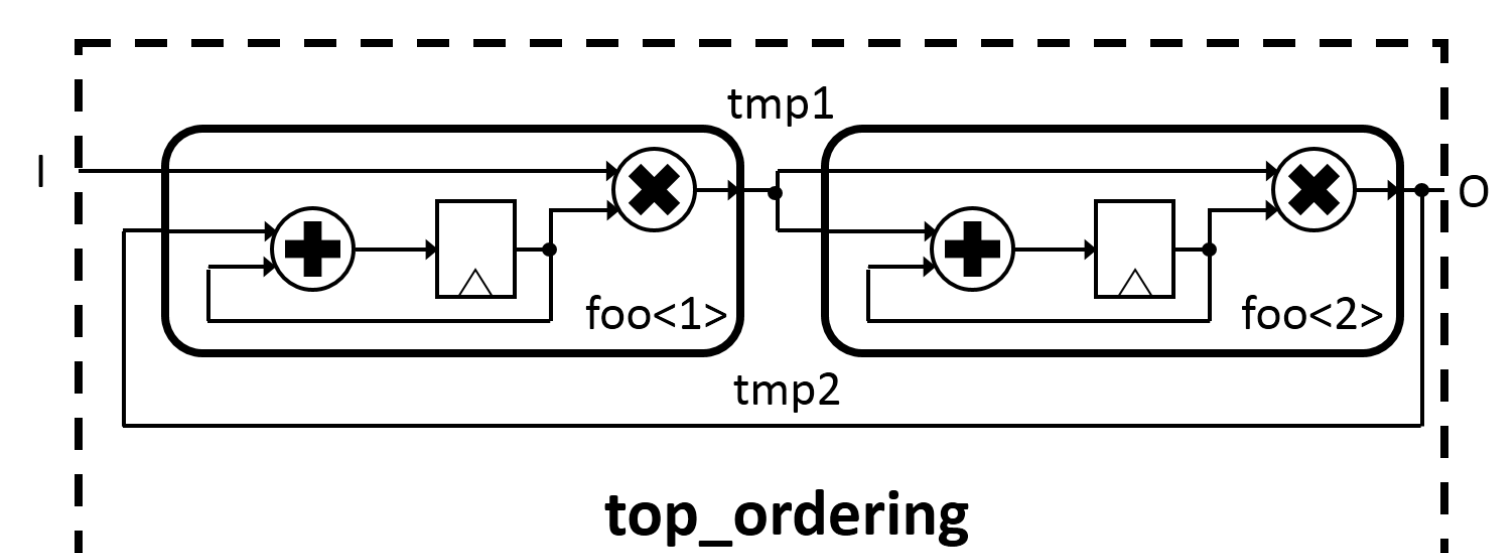
What didn't work so well

Sequential C Language => Concurrent Hardware?

```
void foo(int I1, int I2, int *O) {
    static int L = INIT_VAL; //latch
    *O = I1 * L; //read current-L
    L = I2 + L; //assign next-L
}
```



```
void fxn_ordering_try(int I, int *O) {
    int tmp1 = 0;
    int tmp2 = 0;
    foo<1>(I, tmp2, &tmp1);
    foo<2>(tmp1, tmp1, &tmp2);
    *O = tmp2;
}
```



- See paper for how to get around this using C++ objects

Tech report and source code available

http://www.ece.cmu.edu/calcm/connect_hls